

## CLAIMS

1. A method of setting multicast transfer routes in a multicast network comprising a plurality of points, the multicast transfer routes connecting a given starting point and a plurality of ending points, the multicast network comprising a multicast transfer apparatus provided to each point, a multicast transfer route computing apparatus that computes the multicast transfer routes, and a multicast transfer route setting apparatus that sets the computed multicast transfer routes, the method comprising the following steps:
  - the multicast transfer apparatus measures a traffic state of each direction in which data flow through each link of the network and requests the multicast transfer route computing apparatus to compute multicast transfer routes by transmitting the measured traffic state;
  - the multicast transfer route computing apparatus computes a shortest route with respect to delay connecting the starting point and the ending points based on the measured traffic state, computes delay from each point on the shortest route at the same time, and stores the computed delay in a recording medium;
  - the multicast transfer route computing apparatus computes a greatest delay in the data flow through the computed shortest route;
  - the multicast transfer route computing apparatus compares the greatest delay with a predefined delay condition, re-defines the delay condition if the greatest delay does not satisfy the delay condition, searches, if a condition that the shortest route satisfies is found, a partial route in the computed shortest route that has two

of the same kind or different kinds of the starting point,  
the ending points, and branching points, as ending nodes  
of the partial route, that has none of the starting point,  
the ending points, and branching points in the middle, and  
5 that incurs the greatest cost, removes the searched  
partial route from the computed shortest route thereby to  
divide the multicast transfer route into two route trees,  
and sets a route computed separately as a complementary  
route that complements the removed route to connect the  
10 two route trees;

the multicast transfer route computing apparatus  
informs the multicast transfer route setting apparatus of  
the result of computation; and

the multicast transfer route setting apparatus  
15 sets the multicast transfer route in accordance with the  
informed result of computation.

2. The method as claimed in claim 1,  
wherein

20 when computing the complementary route, the  
multicast transfer route computing apparatus computes a  
route that incurs delay between the starting point and the  
ending points, that is less than an upper limit, and that  
incurs the least cost.

25

3. The method as claimed in claim 2,  
wherein

when computing the complementary route, the  
multicast transfer route computing apparatus selects the  
30 starting point of the complementary route to be computed  
from nodes in the route tree that include the starting  
point, and computes a route of which ending point is the  
ending point of the removed route.

4. The method as claimed in claim 3,  
wherein  
when computing the complementary route, the  
5 multicast transfer route computing apparatus applies a  
known k-th shortest path algorithm to delay, and repeats  
the applying of the k-th shortest path algorithm while the  
delay condition is satisfied.
- 10 5. The method as claimed in claim 4,  
wherein  
when computing the complementary route, the  
multicast transfer route computing apparatus applies the  
known k-th shortest path algorithm to cost, and repeats  
15 the applying of the k-th shortest path algorithm until a  
route that satisfies the delay condition is found.
6. The method as claimed in claim 4,  
wherein  
20 when computing the complementary route, the  
multicast transfer route computing apparatus uses the  
delay from the starting point of the multicast transfer  
route to the starting point of the complementary route  
stored in the recording medium when the shortest route is  
25 computed, and the delay from the ending point of the  
complementary route to each of the ending points of the  
multicast transfer route in order to determine whether the  
result of computation satisfies the delay condition.
- 30 7. An apparatus for computing a multicast  
transfer route in a multicast network, comprising:  
a measurement result receiving unit that  
receives the result of measurement of traffic state in the

multicast network;

- a measurement information storing unit that stores the received result of measurement;
- a measurement result storing unit that causes  
5 the measurement information storing unit to store the result of measurement; and
- a route computing unit that reads the result of measurement from the measurement information storing unit, and computes the multicast transfer route based on the  
10 result of measurement,
  - wherein
  - the route computing unit further comprises:
    - a shortest route delay computing unit that  
15 computes a shortest route with respect to delay connecting the starting point and the ending points based on the measured traffic state, computes delay from each point on the shortest route at the same time, and stores the computed delay in a recording medium;
    - a maximum delay computing unit that computes a  
20 greatest delay in the data flow through the computed shortest route;
    - a maximum cost route searching unit that  
compares the greatest delay with a predefined delay condition, re-defines the delay condition if the greatest  
25 delay does not satisfy the delay condition, searches, if a condition that the shortest route satisfies is found, a partial route in the computed shortest route that has two of the same kind or different kinds of the starting point, the ending points, and branching points as ending nodes of  
30 the partial route, that has none of the starting point, the ending points, and branching points in the middle, and that incurs the greatest cost;
    - a route tree dividing unit that removes the

searched partial route from the computed shortest route thereby to divide the multicast transfer route into two route trees; and

5 a complementary route computing unit that sets a route computed separately as a complementary route that complements the removed route to connect the two route trees.

8. The apparatus as claimed in claim 7,  
10 wherein  
the complementary route computing unit computes a route that incurs delay between the starting point and the ending points, that is less than an upper limit, and that incurs the least cost.

15 9. The apparatus as claimed in claim 8,  
wherein  
the complementary route computing unit selects the starting point of the complementary route to be  
20 computed from nodes in the route tree that includes the starting point, and computes a route of which ending point is the ending point of the removed route.

10. The apparatus as claimed in claim 9,  
25 wherein  
the complementary route computing unit applies a known k-th shortest path algorithm to delay, and repeats the applying of the k-th shortest path algorithm while the delay condition is satisfied.

30 11. The apparatus as claimed in claim 10,  
wherein  
the complementary route computing unit applies

the known k-th shortest path algorithm to cost, and repeats the applying of the k-th shortest path algorithm until a route that satisfies the delay condition is found.

5                   12. The apparatus as claimed in claim 10,  
                  wherein

                  the complementary route computing unit uses the  
                  delay from the starting point of the multicast transfer  
                  route to the starting point of the complementary route  
10                  stored in the recording medium when the shortest route is  
                  computed, and the delay from the ending point of the  
                  complementary route to each of the ending points of the  
                  multicast transfer route in order to determine whether the  
                  result of computation satisfies the delay condition.

15                   13. The apparatus as claimed in claim 7,  
                  further comprising:

                  an indicating unit that indicates the result of  
                  computation by the route computing unit in a control  
20                  message for setting the transfer route; and

                  a transmission unit that transmits the control  
                  message through the multicast transfer route indicated by  
                  the result of computation.

25                   14. The apparatus as claimed in claim 7,  
                  further comprising:

                  a receiving unit that receives a request for  
                  computing the multicast transfer route from a multicast  
                  transfer route setting apparatus; and

30                  a transmitting unit that transmits the result of  
                  computation to the multicast transfer route setting  
                  apparatus.

15. A computer program that causes a computer to compute a multicast transfer route based on the result of measurement of traffic state incurred in links in a multicast network, the computer program comprising the

5 steps of:

- computing the shortest route with respect to delay connecting the starting point and the ending points based on the measured traffic state;
- computing delay from each node on the shortest

10 route at the same time;

- storing the computed delay in a recording medium;
- computing the greatest delay in data flow through the computed shortest route;

15 comparing the greatest delay with a predefined delay condition,

- re-defining, if the greatest delay does not satisfy the delay condition, the delay condition;
- searching, if a condition that the shortest

20 route satisfies is found, a partial route in the computed shortest route that has two of the same kind or different kinds of the starting node, the ending nodes, and branching nodes as ending nodes thereof, that has none of the starting node, the ending nodes, and branching nodes

25 in the middle, and that incurs the greatest cost;

- removing the searched partial route from the computed shortest route thereby to divide the multicast transfer route into two route trees;
- setting a route computed separately as a

30 complementary route that complements the removed route to connect the two route trees.

16. The computer program as claimed in claim 15,

wherein  
when computing the complementary route, the  
computer computes a route that incurs delay between the  
starting node and the ending nodes is less than an upper  
5 limit, and that incurs the least cost.

17. The computer program as claimed in claim 16,  
wherein  
the computer selects the starting node of the  
10 complementary route to be computed from the nodes in the  
route tree that includes the starting node, and computes a  
route of which ending node is the ending node of the  
removed route.

15 18. The computer program as claimed in claim 17,  
wherein  
the computer applies a known k-th shortest path  
algorithm to the delay, and repeats the application while  
the delay condition is satisfied.

20 19. The computer program as claimed in claim 18,  
wherein  
the computer applies the known k-th shortest  
path algorithm to the cost, and repeats the application  
25 until a route that satisfies the delay condition is found.

20. The computer program as claimed in claim 18,  
wherein  
the computer uses the delay from the starting  
30 node to the starting node of the complementary route  
stored in the recording medium when the shortest route is  
computed, and the delay from the ending node of the  
complementary route to the ending nodes in the downstream



thereof in order to determine whether the result of computation satisfies the delay condition.

21. A recording medium storing a computer  
5 program that causes a computer to compute a multicast transfer route based on the result of measurement of traffic state incurred in links in a multicast network,  
wherein  
the recording medium stores the computer program  
10 as claimed in claim 15.

22. A method of multicast label switching in which label switching routes are established for multicast distribution from a multicast source node to a group of  
15 multicast leaf nodes, the method comprising the steps of:  
establishing a point-to-multipoint label switching route of a most upper layer from the multicast source node to all multicast leaf nodes;  
establishing a plurality of label switching  
20 routes of a second layer that configure partial trees of a label switching route of a first layer using second layer labels for respective subgroups of leaf nodes, the subgroup of leaf nodes being extracted as destinations from the group of leaf nodes for which the point-to-  
25 multipoint label switching route has been established;  
allocating traffics addressed to a destination leaf group corresponding to the second layer labels to a corresponding hierarchical label using the first layer label switching route and the second layer label switching  
30 routes by an input label edge router;  
label-switching packets in accordance with a label pair of the first layer and the second layer by a relay label switch router;

if a relay node is designated as a branching node of the point-to-multipoint label switching route, replacing input label pair with output labels corresponding to a plurality of output branches and  
5 copying the input label pair for each output branch;  
switching the input hierarchical label packets to an output line by an output label edge router, and identifying the group of the input hierarchical labels and removing the labels; and  
10 label-switching traffic of each second layer subgroup using point-to-multipoint LSP of the second layer forming the first layer partial tree of a plurality of second layer forming different destination subgroups of the first layer leaf group nodes in the point-to-  
15 multipoint LSP with the first layer label switching route shared.

23. The method as claimed in claim 22,  
wherein  
20 a plurality of label switching routes of a third layer is provided in the multicast label switching routes of the second layer using the sub-tree forming a partial topology of the second layer label switching routes to the leaf nodes configuring a subset of leaf nodes forming the  
25 second layer label switching route;  
if the label switching routes need to be classified into subgroups, establishing label switching routes of a lower layer in a recursive manner;  
performing multicast-label-switching by the  
30 subgroup using the recursively established hierarchical label switching routes.

24. The method as claimed in claim 22,

wherein  
connecting all provider edge routers of a  
provider network accommodating the VPN sites with first  
layer point-to-multipoint multicast LSP in a full-mesh  
5 manner;  
establishing second layer multicast label switch  
routes for each VPN site accommodated in the provider  
network;  
in the case in which the second layer label  
10 switch routes are established, if provider edge routers  
forming the VPN is the leaf node of the multicast label  
switch routes, adjusting the second layer label switching  
route depending on the VPN sites accommodated in each  
leaf;  
15 configuring the second layer label switch route  
in the first layer multicast tree connecting the provider  
edge routers of the VPN.

25. The method as claimed in claim 24,  
20 wherein  
in the case of a multicast distribution route  
having a plurality of different site destinations in the  
VPN site, a third layer multicast distribution route is  
established as the partial tree route of the second layer  
25 multicast distribution route in the third layer under the  
second layer, of which destination leaf nodes are only VPN  
sites corresponding to respective multicast distribution  
routes;  
even the multicast traffic belonging to the same  
30 VPN is distributed to only the VPN sites that wishes to  
receive multicast traffic.

26. The method as claimed in claim 22,

wherein

communication method is provided as a label  
switch router function; and

the communication method is operated as the  
5 input multicast label switch router, the relay multicast  
label switch router, and the output multicast label switch  
router.

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